# Combinatorics and Graph Theory III Tutorial 1 

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## 1 Bipartite graphs

## Definition 1.1

Let $k \in \mathbb{N}$. A graph $G$ is $k$-partite if you can partition its vertices into $k$ disjoint independent sets.
Determining whether a graph is $k$-partite is hard, but in the case $k=2$, there is a simple characterization:

## Theorem 1.2

A graph is bipartite if and only if it has no odd induced cycle.
Two remarks:

- a graph has an odd cycle if and only if it has an odd induced cycle,
- a closed walk is a cycle allowed to reuse an edge/a vertex multiple times. A graph has an odd cycle if and only if it has an odd closed walk.

1. Prove that if a graph is bipartite, then it has no odd induced cycle.
2. Prove that if a graph has no odd induced cycle, then it is bipartite.
hint: If a graph is bipartite, fix a vertex $v$ : there is a simple characterization of each part of the bipartition.

## 2 Minors

A graph $H$ is a minor of a graph $G$ if it can be obtained from $G$ by repeated applications of the following rules:

1. delete a vertex
2. delete an edge
3. contract an edge

If only rules (1) and (2) are used, $H$ is a subgraph of $G$. If only rule (1) is used, $H$ is an induced subgraph of $G$. If only rule (3) is applied and only whenever one endpoint of the edge has degree $2, H$ is a subdivision of $G$.

1. Prove that we can always delete vertices, then delete edges, then contract edges.
2. Prove that forests are exactly the graphs that do not contain $C_{3}$ as a minor.
3. Is $K_{3,3}$ a minor of the Petersen graph ?


## 3 Kuratowski's theorem

A graph is planar if it admits a planar embedding. A planar embedding is a is a drawing of a graph in which edges do not cross outside of vertices. But it is cumbersome to define it properly:

## Definition 3.1

A planar embedding of a graph $G=(V, E)$ is an injective function $\psi: V \rightarrow \mathbb{R}^{2}$ and a set of continuous
functions $\varphi_{u v}:[0,1] \rightarrow R$ such that:

- $\varphi_{u v}(0)=\psi(u)$,
- $\varphi_{u v}(1)=\psi(v)$,
- $\left.\forall t, t^{\prime} \in\right] 0,1\left[, \forall e, e^{\prime} \in E, \varphi_{e}(t)=\varphi_{e}^{\prime}\left(t^{\prime}\right)\right.$ if and only if $e=e^{\prime}$ and $t=t^{\prime}$.

Fortunately, there exists a nice structural characterization of planar graphs:
Theorem 3.2 (Kuratowski's theorem)
A graph is planar if and only if it does not contain any subdivision of $K_{3,3}$ or $K_{5}$.
This also exists in a flavour using minors instead of subdivisions:
Theorem 3.3 (Wagner's theorem)
A graph is planar if and only if it does not contain $K_{3,3}$ or $K_{5}$ as a minor.

1. Using Euler's formula, prove that $K_{3,3}$ and $K_{5}$ are not minors.
2. Argue that if a graph is planar, so are all its minors.
3. Prove that Kuratowski's theorem implies Wagner's theorem.
4. Prove that containing $K_{3,3}$ as a minor is equivalent to having it as a subdivision.
5. Prove that containing $K_{5}$ as a minor implies containing $K_{5}$ or $K_{3,3}$ as a subdivision.

## 4 Surfaces

A surface is a two-dimensional manifold, that is a two-dimensional object that locally looks like the Euclidean plane. Some examples of surfaces: the plane, the projective plane, the sphere, the torus, the Klein bottle.

1. Argue that a graph can be embedded on a sphere if and only if it can be embedded in the plane.
2. does Pac-Man live on a sphere?
3. draw an embedding of $K_{5}$ on a torus.
4. argue that every minor of a graph which can be embedded on a surface can be embedded in this surface.

## 5 Chordal Graphs

## Definition 5.1

A graph is chordal if its only induced cycles have length 3.
An important fact about chordal graphs is that they contain a simplicial vertex:

## Theorem 5.2

Every chordal graph contains a vertex whose neighbourhood is a clique. Such a vertex is called simplicial.

1. What can we say about the induced subgraphs of a chordal graph?

Thus, we can always peel off a chordal graph by repeatedly removing simplicial vertices. But does the reverse process always produce a chordal graph?
2. Let $G$ be a chordal graph, and $K$ a clique of $G$. Add to a $G$ a vertex that sees exactly $K$. Prove that the obtained graph is chordal.

This proves that chordal graphs are exactly those that can be obtained by repeatedly adding simplicial vertices.

